# Implementing a Bootloader for the PIC16F87X 

Author: Mike Garbutt Microchip Technology Inc.

## INTRODUCTION

The PIC16F87X family of microcontrollers has the ability to write to their own program memory. This feature allows a small bootloader program to receive and write new firmware into memory. This application note explains how this can be implemented and discusses the features that may be desirable.
In its most simple form, the bootloader starts the user code running, unless it finds that new firmware should be downloaded. If there is new firmware to be downloaded, it gets the data and writes it into program memory. There are many variations and additional features that can be added to improve reliability and simplify the use of the bootloader, some of which are discussed in this application note.
The general operation of a bootloader is discussed in the OPERATION section. Appendix A contains assembly code for a bootloader developed for the PIC16F877 and key aspects of this bootloader are described in the IMPLEMENTATION section.
For the purpose of this application note, the term "boot code" refers to the bootloader code that remains permanently in the microcontroller and the term "user code" refers to the user's firmware written into FLASH memory by the boot code.

## FEATURES

The more common features a bootloader may have are listed below:

- Code at the Reset location.
- Code elsewhere in a small area of memory.
- Checks to see if the user wants new user code to be loaded.
- Starts execution of the user code if no new user code is to be loaded.
- Receives new user code via a communication channel if code is to be loaded.
- Programs the new user code into memory.


## OPERATION

The boot code begins by checking to see if there is new user code to be downloaded. If not, it starts running the existing user code. If there is new user code to be downloaded, the boot code receives and writes the data into program memory. There are many ways that this can be done, as well as many ways to ensure reliability and ease of use.

## Integrating User Code and Boot Code

The boot code almost always uses the Reset location and some additional program memory. It is a simple piece of code that does not need to use interrupts; therefore, the user code can use the normal interrupt vector at $0 \times 0004$. The boot code must avoid using the interrupt vector, so it should have a program branch in the address range $0 \times 0000$ to $0 \times 0003$.
The boot code must be programmed into memory using conventional programming techniques, and the configuration bits must be programmed at this time. The boot code is unable to access the configuration bits, since they are not mapped into the program memory space. Setting the configuration bits is discussed in the next section.
In order for the boot code to begin executing the user code, it must know where the code starts. Since the boot code starts at the Reset vector, the user code cannot start at this location. There are two methods for placing the starting point of the user code.
One method is to use an ORG directive to force the user code to start at a known location, other than the Reset vector. To start executing the user code, the boot code must branch to this fixed location, and the user code must always use this same location as its start address.
An alternative method is to start the user code at the normal Reset vector and require that the user code has a goto instruction in the first four instructions to avoid the interrupt vector. These four instructions can then be relocated by the boot code and programmed into the area of program memory used by the boot code. This simplifies the development of code for use with the bootloader, since the user code will run when programmed directly into the chip without the boot code present. The boot code must take care of paging and banking so the normal Reset conditions apply before executing the relocated code.

FIGURE 1: INTEGRATING USER CODE WITH BOOT CODE


## Configuration Bits

The configuration bits cannot be changed by the boot code since they are not mapped into the program memory space. This means that the following configuration options must be set at the time that the boot code is programmed into the device and cannot be changed:

## CPx Program Memory Code Protection Enable <br> DEBUG In-Circuit Debugger Mode Enable <br> WRT Program Memory Write Enable <br> CPD Data EEPROM Code Protection Enable <br> LVP Low Voltage In-Circuit Programming Enable <br> BODEN Brown-out Reset Enable <br> PWRTE Power-up Timer Enable <br> WDTE Watchdog Timer Enable <br> FOSCx Oscillator Selection

Most of these configuration options are hardware or design-dependent, and being unable to change them when the user code changes is of no consequence.

The various PIC16F87X devices have different code protection implementations. Please consult the appropriate data sheet for details.

Some devices (such as the PIC16F877), can code protect part of the program memory and prevent internal
writes to this protected section of memory. This can be used to protect the boot code from being overwritten, but also prevents the user code from being code protected, however.

On some devices, code protecting all the program memory still allows internal program memory write cycles. This provides security against the user code being read out of the chip, but does not allow the boot code to be protected from being overwritten.
Data EEPROM Code Protection Enable would normally not need to be set, unless data is programmed into the data EEPROM when the boot code is originally programmed and this data needs to be protected from being overwritten by the user code.

Program Memory Write Enable must be enabled for the boot code to work, since it writes to program memory. Low Voltage In-Circuit Serial Programming (ICSP ${ }^{\text {TM }}$ ) enable only needs to be set if the user wishes to program the PICmicro MCU in-circuit, using logic level signals on the RB3, RB6 and RB7 pins. Since the purpose of the boot code is to program user code into the PICmicro MCU, in most cases, it would be redundant to have facilities for low voltage ICSP.
If the Watchdog Timer is enabled, then the boot code must be written to support the Watchdog Timer and all user code will have to support the Watchdog Timer.

## Determining Whether to Load New Code or to Execute User Code

After a Reset, the boot code must determine whether to download new user code. If no download is required, the bootcode must start execution of existing user code, if available.

There are many ways to indicate whether or not new user code should be downloaded. For example, by testing a jumper or switch on a port pin, polling the serial port for a particular character sequence, or reading an address on the $\mathrm{I}^{2} \mathrm{C}^{\text {TM }}$ bus. The particular method chosen depends on the way that user code is transferred into the microcontroller. For example, if the new user code is stored on an $I^{2} C$ EEPROM that is placed in a socket on the board, then an address in the EEPROM could be read to determine whether a new EEPROM is present.

If an error occurred while downloading new user code, or the bootloader is being used for the first time, there might not be valid user code programmed into the microcontroller. The boot code should not allow faulty user code to start executing, because unpredictable results could occur.

## Receiving New User Code to Load into Program Memory

There are many ways that the microcontroller can receive the new firmware to be written into program memory. A few examples are from a PC over a serial port, from a serial EEPROM over an $\mathrm{I}^{2} \mathrm{C}$ or $\mathrm{SPI}{ }^{T M}$ bus, or from another microcontroller through the parallel slave port.

The boot code must be able to control the reception of data, since it cannot process any data sent to it while it is writing to its own program memory. In the case of data being received via RS-232, there must be some form of flow control to avoid data loss.
The data received by the boot code will usually contain more than just program memory data. It will normally contain the address to which the data is to be written and perhaps a checksum to detect errors. The boot code must decode, verify and store the data, before writing it into program memory. The available RAM (GPR registers) of the device limits the amount of data that can be received before writing it to program memory.

## Programming the FLASH Program Memory

The PIC16F87X devices have special function registers that are used to write data to program memory. There is a specific sequence of writes to these registers that must be followed to reduce the chances of an unintended program memory write cycle occurring. Because code cannot be executed from the FLASH program memory while it is being written, program execution halts for the duration of the write cycle. Program memory is written one word at a time.

## Error Handling

There are several things that can go wrong during execution of the boot code or user code. The bootloader should handle the following error conditions:

- No valid user code written into the chip.
- Error in incoming data.
- Received user code does not have any code at its Reset vector.
- Received user code overlaps boot code.
- User code causes execution into the boot code area.
If the bootloader is being used for the first time, or if the user code is partially programmed because of a previous error, there might not be valid user code programmed into the microcontroller. The boot code should not allow potentially faulty user code to start executing.
The transfer of data can be interrupted, which will cause the boot code to stop receiving data. There are several ways to handle this depending on how the data is being received. For example, the boot code may be able to time-out and request the data to be sent again. The simplest method is to wait, trying to receive more data with no time-out, until the user intervenes and resets the device. Since the boot code needs to leave the most possible program memory space for the user code and also be reliable, the smallest, simplest implementation is often the best.
Incoming data may be corrupted by noise or some other temporary interruption, and this should be detected, otherwise, incorrect data could be programmed. A checksum or other error detection method can be used.

Incorrect use of flow control can result in data being sent to the PICmicro MCU while it is not ready to receive data. This can cause overrun errors that should be handled by the boot code. Once an overrun has occurred, the data is lost and this is essentially the same as a data transfer interruption, discussed above.
In some cases, data could be sent to the microcontroller before the boot code is running, causing part of the data to be lost. If this type of error is possible, then it should be detected. This error may manifest itself as user code that does not seem to have any code at the Reset location and can be detected by checking the addresses being programmed. An alternative is to generate a checksum on all the code that is written into program memory and transmit this to the user for verification, after programming has been completed.

The code developer should take care that the user code does not use the same program memory space that the boot code uses. The exception is the user code at the Reset location that can be relocated, as explained earlier. If the user code does try to use program memory that contains boot code, the boot code should detect the conflicting address and not overwrite itself. In some devices, part of the program memory can be code protected to prevent internal writes to the part of the memory that contains the main boot code. Note that this does not apply to all PIC16F87X devices.
Faulty user code, or a brown-out condition that corrupts the program counter, can cause execution to jump to an unprogrammed memory location and possibly run into the start of the boot code. If the user code at the Reset location is being relocated, as explained earlier, then execution can enter the boot code area if a program branch does not occur in these four relocated instructions. The boot code should trap the program execution to avoid these errors from causing any unintended operation.
When an error is detected, it is useful to indicate this in some way. This can be as simple as turning on an LED, or sending a byte out the serial port. If the system includes a display and the display drivers are incorporated into the boot code, then more sophisticated error messages can be used.

FIGURE 2: FLOWCHART FOR BOOTLOADER


FIGURE 3: SCHEMATIC SHOWING SERIAL PORT AND TEST PIN


## IMPLEMENTATION

## How this Bootloader Works

The boot code in Appendix A implements a bootloader in a PIC16F877 device. It uses the USART to receive data with hardware handshaking, tests a pin to decide if new user code should be received and includes many of the features discussed in this application note.

## Integrating User Code and Boot Code

The code at the Reset location (ResetVector) writes to PCLATH. To set the page bits, it then jumps to the rest of the boot code in upper memory. The main code is in the upper 224 bytes of memory starting at address 0x1F20 (StartOfBoot). The first instructions at this location trap accidental entry into the boot code. The main bootloader routine starts at the address labeled Main.
The boot code requires that the user code includes a goto instruction in the first four locations after the Reset vector and relocates these four instructions into the boot code section (StartUserCode). This simplifies the development of code for use with the bootloader, since the same user code will also run when programmed directly into the chip, without the boot code present. The boot code changes to bank 0 and clears PCLATH before executing the relocated code, so that the normal Reset conditions apply. If a program branch does not occur in the four relocated instructions, then program execution is trapped in an endless loop to avoid any unintended operation.
The boot code must be programmed into the PIC16F877 using conventional programming techniques and the configuration bits are programmed at the same time. The configuration bits are defined with a __CONFIG directive and cannot be accessed by the boot code, because they are not mapped into the program memory space. The boot code does not use a Watchdog Timer.

## Determining Whether to Load new Code or to Execute User Code

The boot code tests port pin RB0 to determine whether new user code should be downloaded. If a download is required, then the boot code branches to the Loader routine that receives the data and writes it into program memory.
If pin RB0 does not indicate that new user code should be loaded, then a program memory location (labeled CodeStatus) is read with routine FlashRead to determine whether there is valid user code in the device. If there is valid user code, the boot code transfers execution to the user code by branching to location StartUserCode. Otherwise, execution is trapped in an endless loop to avoid this error from causing any unintended operation.

## Receiving New User Code to Load into Program Memory

The boot code receives the new firmware as a standard Intel ${ }^{\circledR}$ hex file (INHX8M format), using the USART in Asynchronous Receiver mode (hex format defined in Appendix B). It is assumed that a PC will be used to send this file via an RS-232 cable, connected to a COM port. Hardware handshaking allows the boot code to stop the PC from transmitting data while FLASH program memory is being written. Since the PICmicro device halts program execution while the FLASH write occurs, it cannot read data from the USART during this time.
Hardware handshaking (described in Appendix C) is implemented using port pin RB1 as the RTS output and RB2 as the CTS input. The USART is set to 8-bit Asynchronous mode at 9600 baud in the Serialsetup routine. The SerialReceive routine enables transmission with the RTS output and waits until a data byte has been received by the USART, before returning with the data. The SerialTransmit routine checks the CTS input until a transmission is allowed and then sends a byte out the USART. This is used for transmitting progress indication data back to the PC.
The boot code receives the hex file, one line at a time and stops transmission after receiving each line, while received data is programmed into program memory.

## Decoding the Hex File

The boot code remains in a loop, waiting until a colon is received. This is the first character of a line of the hex file. The following four pairs of characters are received and converted into bytes, by calling the GetHexByte routine. The number of bytes (divided by two to get the number of words) and the address (divided by two to get a word address) are saved, and the record type is checked for a data record, or end of file record.
If the record type shows that the line contains program memory data, then this data is received, two pairs of characters at a time (using the GetHexByte routine), and is stored in an array. The checksum at the end of the line is received and checked, to verify that there were not any errors in the line.
Once the hex file line has been received, hardware handshaking is used to stop further transmission, while the data is written into the program memory. The <CR> and <LF> characters that are sent at the end of the line are ignored. This gives the handshaking time to take effect by ignoring the byte being transmitted, when the handshaking signal is asserted. Once the data from the line has been programmed, the following lines are received and programmed in the same way, until the line indicating the end of the file has been received. A success indication ' $S$ ' is then transmitted out the USART (by the FileEnd routine) and the boot code waits for a Reset.

## Programming the FLASH Program Memory

Data is written to the FLASH program memory using special function registers. The address is written to the EEADR and EEADRH registers and the first two bytes of data are written to EEDATA and EEDATH. The FlashWrite routine is then executed, which writes the data into program memory. The address is then incremented and the next two data bytes are written. This is repeated until all the data from the line of the hex file has been programmed into the FLASH program memory.

## Error Handling

There are several things that can go wrong during execution of the boot code or user code, and a number of these error conditions are handled by the boot code. If an error occurs, the boot code traps it by executing an infinite loop, until the user intervenes and resets the device. If an error is detected in the incoming data, then a failure indication ' $F$ ' is transmitted. This does not occur in the case of an overflow error, or if the data transmission is halted.
If the bootloader is being used for the first time, or if the user code is partially programmed because of a previous error, there might not be valid user code programmed into the microcontroller. The boot code handles this by writing a status word ( $0 \times 3 \mathrm{fff}$ ) at a location labeled CodeStatus, before programming the FLASH device, and then writing a different status word ( $0 \times 0000$ ) to this same location, when programming of the user code has been completed. The boot code tests this location and only starts execution of the user code, if it sees that the user code was successfully programmed. When the boot code is originally programmed into the PICmicro MCU, the status word indicates that there is not valid user code in the device.

The transfer of data can be interrupted. In this case, the boot code waits, trying to receive more data with no time-out, until the user intervenes and resets the device. Noise, or a temporary interruption, may corrupt incoming data. The Intel hex file includes a checksum on each line and the boot code checks the validity of each line by verifying the checksum.

Incorrect use of flow control can result in data being sent to the PIC16F877, while it is not ready to receive data. This can cause an overrun error in the USART. Once an overrun has occurred, the USART will not move any new data into the receive FIFO and the boot code will be stuck in a loop waiting for more data. This effectively traps the error until the user intervenes by resetting the device.
If the user starts transmitting a hex file before the boot code is running, the boot code may miss the first lines of the file. Since all the lines of a hex file have the same format, it is not normally possible to determine whether the line being received is the first line of the hex file. However, since MPASM generates hex files with addresses in ascending order, the first valid line of the
hex file should contain the code for the Reset vector which is checked by the boot code.

The user code may try to use program memory locations that contain boot code. This is detected by checking the address being programmed and detecting conflicting addresses. The boot code will not overwrite itself and is not code protected.
Faulty user code, or noise that corrupts the program counter, can cause execution to jump to an unprogrammed memory location and possibly run into the start of the boot code. The first instructions in the boot code are an infinite loop that traps execution into the boot code area.
Because the first four instructions in program memory are relocated in the boot code implementation, there must be a program branch within these four instructions. If there is no program branch, then execution is trapped by the boot code.

## Using the Bootloader

The procedure for using the bootloader is as follows:

- On the PC, set up the serial port baud rate and flow control (hardware handshaking).
- Connect the serial port of the PIC16F87X device to the serial port of the PC.
- Press the switch to pull pin RB0 low.
- Power up the board to start the boot code running.
- The switch on RBO can be released if desired.
- From the PC, send the hex file to the serial port.
- A period '.' will be received from the serial port for each line of the hex file that is sent.
- An 'S' or 'F' will be received to indicate success or failure.
- The user must handle a failure by resetting the board and starting over.
- Release the switch to set pin RB0 high.
- Power-down the board and power it up to start the user code running.
On the PC, there are several ways to set up the serial port and to transfer data. This also differs between operating systems.
A terminal program allows the user to set up and send data to a serial port. In most terminal programs, an ASCII or text file can be sent and this option should be used to send the hex file. A terminal program will also show data received on the serial port and this allows the user to see the progress '.' indicators and the success ' $S$ ' or failure ' $F$ ' indicators. There are many terminal programs available, some of which are available free on the Internet. This boot code was tested using Tera Term Pro, Version 2.3. The user should be aware that some popular terminal programs contain bugs.

A serial port can be set up in a DOS window, using the MODE command and a file can be copied to a serial port, using the COPY command. When using Windows ${ }^{\circledR} 95 / 98$, the MODE command does not allow the handshaking signals to be configured. This makes it difficult to use the COM port in DOS. When using Windows $\mathrm{NT}^{\circledR}$ or Windows $2000^{\circledR}$, the following commands can be used to send a hex file named filename. hex to serial port COM1:

MODE COM1: BAUD=9600 PARITY=N DATA=8 STOP=1 to=off $x o n=o f f \quad o d s r=o f f$ octs=on dtr=off rts=on idsr=off

COPY filename.hex COM1:

## Resources Used

The boot code coexists with the user code on the PIC16F877 and many of the resources used by the boot code can also be used by the user code. The boot code uses the resources listed in Table 1.

## TABLE 1: RESOURCES USED BY THE BOOT CODE

| Resource | Amount |
| :--- | :--- |
| Program memory | 224 words |
| Data memory | 72 bytes |
| I/O pins | 5 pins |
| Peripherals | USART |

The program memory used by the boot code cannot be used for user code, although it is possible to call some of the subroutines implemented in the boot code to save space. The user code can use all the data memory.
The USART can be used by the user code with the two I/O pins for the USART and the I/O pins used for handshaking. The I/O pin used to indicate that the boot code should load new user code, is connected to a switch or jumper. This can be isolated with a resistor and used as an output, so that it is possible to use all the I/O pins used by the bootloader.
In summary, all resources used by the boot code, except program memory, can also be used by the user code.

## CONCLUSION

Using a bootloader is an efficient way to allow firmware upgrades in the field. Less than $3 \%$ of the total program memory is used by the boot code and the entire program memory available on a PIC16F877 can be programmed in less than one minute at 19,200 baud.
The cost of fixing code bugs can be reduced with a bootloader. Products can be upgraded with new features in the field, adding value and flexibility to the products. The ability to upgrade in the field is an added feature and can enhance the value of a product.


[^0]



|  |  | $\begin{array}{l\|l} \text { M } \\ \text { 㠭 } \\ \text { M } \end{array}$ |  | $\begin{aligned} & U \\ & \infty \\ & \infty \\ & \infty \\ & 0 \end{aligned}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\begin{aligned} & \text { m } \\ & \underset{\sim}{\sim} \\ & \text { I } \end{aligned}$ |  | $\begin{gathered} \text { H } \\ \underset{y}{\mid} \\ \text { H} \end{gathered}$ |  | $\begin{aligned} & \mathbb{H} \\ & \underset{\sim}{\mid} \\ & \underset{H}{2} \end{aligned}$ |  |


| 00159 |  | movwf | EEDATA |  |
| :---: | :---: | :---: | :---: | :---: |
| 00160 |  | call | FlashWrite | ；write new CodeStatus word |
| 00161 |  |  |  |  |
| 00162 |  | call | SerialSetup | ；set up serial port |
| 00163 |  |  |  |  |
| 00164 |  |  |  |  |
| 00165 ；Get new line of hex file starting with＇： |  |  |  |  |
| 00166 ；Get first 8 bytes after＇：＇and extract address and number of bytes |  |  |  |  |
| 00167 |  |  |  |  |
| 00168 | GetNewLine： |  | call | SerialReceive | ；get new byte from serial port |
| 00169 |  | xorlw | ＇：＇ | ；check if＇：＇received |
| 00170 |  | btfss | STATUS，Z |  |
| 00171 |  | goto | GetNewLine | ；if not then wait for next byte |
| 00172 |  |  |  |  |
| 00173 |  | clrf | Checksum | ；start with checksum zero |
| 00174 |  |  |  |  |
| 00175 |  | call | GetHexByte | ；get number of program data bytes in line |
| 00176 |  | andlw | 0x1F | ；limit number in case of error in file |
| 00177 |  | movwf | NumWords |  |
| 00178 |  | bcf | StATUS，C |  |
| 00179 |  | rrf | NumWords，F | ；divide by 2 to get number of words |
| 00180 |  |  |  |  |
| 00181 |  | call | GetHexByte | ；get upper half of program start address |
| 00182 |  | movwf | AddressH |  |
| 00183 |  |  |  |  |
| 00184 |  | call | GetHexByte | ；get lower half of program start address |
| 00185 |  | movwf | AddressL |  |
| 00186 |  |  |  |  |
| 00187 |  | bcf | STATUS，C |  |
| 00188 |  | rrf | Addressh，F | ；divide address by 2 to get word address |
| 00189 |  | rrf | AddressL，F |  |
| 00190 |  |  |  |  |
| 00191 |  | call | GetHexByte | ；get record type |
| 00192 |  | xorlw | 0x01 |  |
| 00193 |  | btfsc | STATUS，Z | ；check if end of file record（0x01） |
| 00194 |  | goto | FileDone | ；if end of file then all done |
| 00195 |  |  |  |  |
| 00196 |  | movf | HexByte，W |  |
| 00197 |  | xorlw | 0x00 |  |
| 00198 |  | btfss | STATUS，Z | ；check if regular line record（0x00） |
| 00199 |  | goto | LineDone | ；if not then ignore line and send＇．＇ |
| 00200 |  |  |  |  |
| 00201 |  | movlw | $0 \mathrm{xe0}$ |  |
| 00202 |  | addwf | Addressh，W | ；check if address＜0x2000 |
| 00203 |  | btfsc | STATUS，C | ；which is ID locations and config bits |
| 00204 |  | goto | LineDone | ；if so then ignore line and send＇．＇ |
| 00205 |  |  |  |  |


| $\begin{array}{ll} U & 爪 \\ \infty & \mathbb{1} \\ 0 & \stackrel{1}{\lambda} \\ 0 \end{array}$ | U N |  | $\begin{aligned} & \underset{\sim}{\sim} \\ & \underset{\sim}{-} \end{aligned}$ |  | $\begin{array}{ll} \text { U } & 0 \\ \\ \text { N } \\ \text { N } \end{array}$ | $\begin{array}{ll} U & \underset{\sim}{C} \\ \\ \underset{\sim}{c} & 0 \end{array}$ |  |  | $$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | － | N M サ に |  | へ $\quad$ の |  |  | $\bigcirc$ | m +1 |  | ¢ U 도 |
| $\stackrel{0}{\text { m }}$ | $\stackrel{7}{4}$ | $\begin{array}{llll}7 \\ 7 & 7 \\ 4 & 4 \\ 4\end{array}$ | $\stackrel{+}{4}$ |  |  |  | ก ก ${ }_{\text {¢ }}$ N |  |  | 邱 |
| 缶 |  | 式嵒受孚 | $\stackrel{\text { H }}{\text {－}}$ |  |  |  | 哣孚舀 |  |  | 孚孚孚孚 |









1FA0



|  | $\begin{aligned} & \text { H } \\ & \text { 世 } \\ & \text { 心 } \end{aligned}$ | MN N N M N N |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\begin{array}{lll} \infty & \text { 氐 } \\ \text { 出畠舀 } \\ \text { - } \end{array}$ |  |  |  |  |




| $\begin{aligned} & 0 \\ & \infty \\ & \underset{\sim}{1} \end{aligned}$ |  |  |  |
| :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { M1 } \\ & \text { U } \\ & \end{aligned}$ |  |  |  |

return


$\infty$
0
0
0

| 6 <br> 0 <br>  <br> + | ＋ | U | 6 | の | $\infty$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | ［1］ | － | ［ | の | $\bigcirc$ |
|  | ［ | ［午 | ［ | － | － |
|  |  |  | $\cdots$ | － | $\bigcirc$ |
|  | 上 | 6 | － | $\infty$ | の |
| ［1］ | ［1］ | ［1 | ［9 | ［1 | ［1］ |
| ［ | ［ | ［ | ［ | ［ | ［ |
|  | － |  | － | $\Gamma$ | $\Gamma$ |



[^1]H
－1
－
－


## APPENDIX B: HEX FILE FORMAT

MPASM generates an 8-bit Intel hex file (INHX8M) by default. The lines of this hex file all have the following format:

```
: BBAAAATTHHHH... . HHCC
```

A colon precedes each line and is followed by hexadecimal digits in ASCII format.
BB is a 2-digit hexadecimal byte count representing the number of data bytes that will appear on the line. This is a number from $0 \times 00$ to $0 \times 10$ and is always even because the PIC16F87X parts have a 14-bit wide memory and use two bytes for every program memory word.
AAAA is a 4-digit hexadecimal address representing the starting byte address of the data bytes that follow. To get the actual program memory word address, the byte address must be divided by two.

TT is a 2-digit hexadecimal record type that indicates the meaning of the data on the line. It is $0 \times 00$ for a regular data record and $0 \times 01$ for an end of file record. The boot code ignores all other record types.

HH are 2-digit hexadecimal data bytes that correspond to addresses, incrementing sequentially from the starting address earlier in the line. These bytes come in low byte, high byte pairs, corresponding to each 14-bit program memory word.
CC is a 2-digit hexadecimal checksum byte, such that the sum of all bytes in the line including the checksum, is a multiple of 256 . The initial colon is ignored.

The code in Example B-1 will generate a line in a hex file as shown in Figure B-1.

EXAMPLE B-1: CODE TO GENERATE A HEX FILE

| ORG | 0x17A |
| :--- | :--- |
| movlw | 0xFF |
| movwf | PORTB |
| bsf | STATUS,RP0 |
| movwf | TRISA |
| clrf | TRISB |
| bcf | STATUS,RP0 |

## FIGURE B-1: LINE OF HEX FILE

: OC02F400FF30860083168500860183120F


Result of addition ${ }^{(1)}$ mod 256 is zero

Second program memory word is $0 \times 0086$
This corresponds to an instruction MOVWF $0 \times 06^{(2)}$

First program memory word is $0 \times 30 \mathrm{FF}$
This corresponds to an instruction MOVLW 0xFF

Record type is $0 \times 00$ indicating a regular data record

Address of first program memory word is $0 \times 02 \mathrm{~F} 4 \div 2=0 \times 017 \mathrm{~A}$

Number of data bytes is $0 x 0 \mathrm{C}$
Number of program memory words is $0 \times 0 C \div 2=0 \times 06$

Note 1: The calculation to test the checksum adds every byte (pair of digits) in the line of the hex file, including the checksum itself.

2: The label PORTB is defined as $0 x 06$ in the standard include file for the PIC16F877.

## APPENDIX C: RS-232 HARDWARE HANDSHAKING SIGNALS

Understanding hardware flow control can be confusing, because of the terminology used and the slightly different way that handshaking is now implemented, compared to the original specification.
RS-232 hardware handshaking was specified in terms of communication between Data Terminal Equipment (DTE) and Data Communications Equipment (DCE). The DTE (e.g., computer terminal) was always faster than the DCE (e.g., modem) and could receive data without interruption. The hardware handshaking protocol required that the DTE would request to send data to the DCE (with the request to send RTS signal) and that the DCE would then indicate to the DTE that it was cleared to send data (with the clear to send CTS signal). Both RTS and CTS were, therefore, used to control data flow from the DTE to the DCE.

The Data Terminal Ready (DTR) signal was defined so that the DTE could indicate to the DCE that it was attached and ready to communicate. The Data Set Ready (DSR) signal was defined to enable the DCE to indicate to the DTE that it was attached and ready to communicate. These are higher level signals not generally used for byte by byte control of data flow, although they can be used for this purpose.
Most RS-232 connections use 9-pin DSUB connectors. A DTE uses a male connector and a DCE uses a female connector. The signal names are always in terms of the DTE, so the RTS pin on the female connector of the DCE is an input and is the RTS signal from the DTE.

Over time, the clear distinction between the DTE and DCE has been lost. In many instances, two DTE devices are connected together. In other cases, the DCE device is able to send data at a rate that is too high for the DTE to receive continuously. In practice, the DTR output of the DTE has come to be used to control the flow of data to the DTE and now indicates that the DCE (or other DTE) may send data. It no longer indicates a request to send data to the DCE.
It is common for a DTE to be connected to another DTE (e.g., two computers), and in this case, they will both have male connectors and the cable between them will have two female connectors. This is known as a null modem cable. The cable is usually wired in such a way that each DTE looks like a DCE to the other DTE. To achieve this, the RTS output of one DTE is connected to the CTS input of the other DTE and vice versa. Each DTE device will use its RTS output to allow the other DTE device to transmit data and will check its CTS input to determine whether it is allowed to transmit data.

FIGURE C-1: DTE TO DCE CONNECTION


Data Terminal Equipment DSUB9 male connector


DTE to DCE RS-232 cable DSUB9 female to DSUB9 male connector


Data Communication Equipment DSUB9 female connector

FIGURE C-2: DTE TO DTE CONNECTION


Data Terminal Equipment DSUB9 male connector


DTE to DTE RS-232 null modem cable DSUB9 female to DSUB9 female connector


Data Terminal Equipment DSUB9 male connector

## Note the following details of the code protection feature on PICmicro ${ }^{\circledR}$ MCUs.

- The PICmicro family meets the specifications contained in the Microchip Data Sheet.
- Microchip believes that its family of PICmicro microcontrollers is one of the most secure products of its kind on the market today, when used in the intended manner and under normal conditions.
- There are dishonest and possibly illegal methods used to breach the code protection feature. All of these methods, to our knowledge, require using the PICmicro microcontroller in a manner outside the operating specifications contained in the data sheet. The person doing so may be engaged in theft of intellectual property.
- Microchip is willing to work with the customer who is concerned about the integrity of their code.
- Neither Microchip nor any other semiconductor manufacturer can guarantee the security of their code. Code protection does not mean that we are guaranteeing the product as "unbreakable".
- Code protection is constantly evolving. We at Microchip are committed to continuously improving the code protection features of our product.
If you have any further questions about this matter, please contact the local sales office nearest to you.

Information contained in this publication regarding device applications and the like is intended through suggestion only and may be superseded by updates. It is your responsibility to ensure that your application meets with your specifications. No representation or warranty is given and no liability is assumed by Microchip Technology Incorporated with respect to the accuracy or use of such information, or infringement of patents or other intellectual property rights arising from such use or otherwise. Use of Microchip's products as critical components in life support systems is not authorized except with express written approval by Microchip. No licenses are conveyed, implicitly or otherwise, under any intellectual property rights.

## Trademarks

The Microchip name and logo, the Microchip logo, FilterLab, KeELoQ, microID, MPLAB, PIC, PICmicro, PICMASTER, PICSTART, PRO MATE, SEEVAL and The Embedded Control Solutions Company are registered trademarks of Microchip Technology Incorporated in the U.S.A. and other countries.
dsPIC, ECONOMONITOR, FanSense, FlexROM, fuzzyLAB, In-Circuit Serial Programming, ICSP, ICEPIC, microPort, Migratable Memory, MPASM, MPLIB, MPLINK, MPSIM, MXDEV, PICC, PICDEM, PICDEM.net, rfPIC, Select Mode and Total Endurance are trademarks of Microchip Technology Incorporated in the U.S.A.

Serialized Quick Turn Programming (SQTP) is a service mark of Microchip Technology Incorporated in the U.S.A.

All other trademarks mentioned herein are property of their respective companies.
© 2002, Microchip Technology Incorporated, Printed in the U.S.A., All Rights Reserved.

3 Printed on recycled paper.


[^2]
## Worldwide Sales and SERVICE

## AMERICAS

## Corporate Office

2355 West Chandler Blvd.
Chandler, AZ 85224-6199
Tel: 480-792-7200 Fax: 480-792-7277
Technical Support: 480-792-7627
Web Address: http://www.microchip.com

## Rocky Mountain

2355 West Chandler Blvd.
Chandler, AZ 85224-6199
Tel: 480-792-7966 Fax: 480-792-7456

## Atlanta

500 Sugar Mill Road, Suite 200B
Atlanta, GA 30350
Tel: 770-640-0034 Fax: 770-640-0307

## Boston

2 Lan Drive, Suite 120
Westford, MA 01886
Tel: 978-692-3848 Fax: 978-692-3821

## Chicago

333 Pierce Road, Suite 180
Itasca, IL 60143
Tel: 630-285-0071 Fax: 630-285-0075

## Dallas

4570 Westgrove Drive, Suite 160
Addison, TX 75001
Tel: 972-818-7423 Fax: 972-818-2924

## Detroit

Tri-Atria Office Building
32255 Northwestern Highway, Suite 190
Farmington Hills, MI 48334
Tel: 248-538-2250 Fax: 248-538-2260

## Kokomo

2767 S. Albright Road
Kokomo, Indiana 46902
Tel: 765-864-8360 Fax: 765-864-8387

## Los Angeles

18201 Von Karman, Suite 1090
Irvine, CA 92612
Tel: 949-263-1888 Fax: 949-263-1338

## New York

150 Motor Parkway, Suite 202
Hauppauge, NY 11788
Tel: 631-273-5305 Fax: 631-273-5335

## San Jose

Microchip Technology Inc.
2107 North First Street, Suite 590
San Jose, CA 95131
Tel: 408-436-7950 Fax: 408-436-7955

## Toronto

6285 Northam Drive, Suite 108
Mississauga, Ontario L4V 1X5, Canada
Tel: 905-673-0699 Fax: 905-673-6509

## ASIA/PACIFIC

## Australia

Microchip Technology Australia Pty Ltd
Suite 22, 41 Rawson Street
Epping 2121, NSW
Australia
Tel: 61-2-9868-6733 Fax: 61-2-9868-6755

## China - Beijing

Microchip Technology Consulting (Shanghai)
Co., Ltd., Beijing Liaison Office
Unit 915
Bei Hai Wan Tai Bldg.
No. 6 Chaoyangmen Beidajie
Beijing, 100027, No. China
Tel: 86-10-85282100 Fax: 86-10-85282104

## China - Chengdu

Microchip Technology Consulting (Shanghai
Co., Ltd., Chengdu Liaison Office
Rm. 2401, 24th Floor,
Ming Xing Financial Tower
No. 88 TIDU Street
Chengdu 610016, China
Tel: 86-28-6766200 Fax: 86-28-6766599

## China - Fuzhou

Microchip Technology Consulting (Shanghai)
Co., Ltd., Fuzhou Liaison Office
Unit 28F, World Trade Plaza
No. 71 Wusi Road
Fuzhou 350001, China
Tel: 86-591-7503506 Fax: 86-591-7503521

## China - Shanghai

Microchip Technology Consulting (Shanghai)
Co., Ltd.
Room 701, Bldg. B
Far East International Plaza
No. 317 Xian Xia Road
Shanghai, 200051
Tel: 86-21-6275-5700 Fax: 86-21-6275-5060

## China - Shenzhen

Microchip Technology Consulting (Shanghai)
Co., Ltd., Shenzhen Liaison Office
Rm. 1315, 13/F, Shenzhen Kerry Centre,

## Renminnan Lu

Shenzhen 518001, China
Tel: 86-755-2350361 Fax: 86-755-2366086

## Hong Kong

Microchip Technology Hongkong Ltd.
Unit 901-6, Tower 2, Metroplaza
223 Hing Fong Road
Kwai Fong, N.T., Hong Kong
Tel: 852-2401-1200 Fax: 852-2401-3431

## India

Microchip Technology Inc
India Liaison Office
Divyasree Chambers
1 Floor, Wing A (A3/A4)
No. 11, O'Shaugnessey Road
Bangalore, 560 025, India
Tel: 91-80-2290061 Fax: 91-80-2290062

## Japan

Microchip Technology Japan K.K.
Benex S-1 6F
3-18-20, Shinyokohama
Kohoku-Ku, Yokohama-shi
Kanagawa, 222-0033, Japan
Tel: 81-45-471-6166 Fax: 81-45-471-6122

## Korea

Microchip Technology Korea
168-1, Youngbo Bldg. 3 Floor
Samsung-Dong, Kangnam-Ku
Seoul, Korea 135-882
Tel: 82-2-554-7200 Fax: 82-2-558-5934

## Singapore

Microchip Technology Singapore Pte Ltd.
200 Middle Road
\#07-02 Prime Centre
Singapore, 188980
Tel: 65-334-8870 Fax: 65-334-8850

## Taiwan

Microchip Technology Taiwan
11F-3, No. 207
Tung Hua North Road
Taipei, 105, Taiwan
Tel: 886-2-2717-7175 Fax: 886-2-2545-0139

## EUROPE

## Denmark

Microchip Technology Nordic ApS
Regus Business Centre
Lautrup hoj 1-3
Ballerup DK-2750 Denmark
Tel: 4544209895 Fax: 4544209910

## France

Microchip Technology SARL
Parc d'Activite du Moulin de Massy
43 Rue du Saule Trapu
Batiment A - ler Etage
91300 Massy, France
Tel: 33-1-69-53-63-20 Fax: 33-1-69-30-90-79

## Germany

Microchip Technology GmbH
Gustav-Heinemann Ring 125
D-81739 Munich, Germany
Tel: 49-89-627-144 0 Fax: 49-89-627-144-44

## Italy

Microchip Technology SRL
Centro Direzionale Colleoni
Palazzo Taurus 1 V. Le Colleoni 1
20041 Agrate Brianza
Milan, Italy
Tel: 39-039-65791-1 Fax: 39-039-6899883

## United Kingdom

Arizona Microchip Technology Ltd.
505 Eskdale Road
Winnersh Triangle
Wokingham
Berkshire, England RG41 5TU
Tel: 441189215869 Fax: 44-118 921-5820


[^0]:    SOURCE CODE - FILE BOOT877.ASM
    APPENDIX A:
    MPASM 02.40 Released
    BOOT877.ASM 6-26-2000 14:58:44
    LINE SOURCE TEXT

    LOC OBJECT CODE
    VALUE

[^1]:    $\begin{array}{ll}\text { r } \\ -1 \\ 0 \\ 0 & 0 \\ m\end{array}$
    $\begin{array}{ll}\circ & 0 \\ \infty & \infty \\ 0 & 0 \\ m & 0\end{array}$
    $U$
    0
    +
    $H$
    $H$

[^2]:    Microchip received QS-9000 quality system certification for its worldwide headquarters, design and wafer fabrication facilities in Chandler and Tempe, Arizona in July 1999. The Company's quality system processes and procedures are QS-9000 compliant for its PICmicro ${ }^{\oplus}$ 8-bit MCUs, KEELOQ ${ }^{\oplus}$ code hopping devices, Serial EEPROMs and microperipheral products. In addition, Microchip's quality system for the design and manufacture of development systems is ISO 9001 certified.

