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Implementing a Bootloader for the PIC16F87X

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INTRODUCTION

The PIC16F87X family of microcontrollers has the ability to write to their own program memory. This feature allows a small bootloader program to receive and write new firmware into memory. This application note explains how this can be implemented and discusses the features that may be desirable.

In its most simple form, the bootloader starts the user code running, unless it finds that new firmware should be downloaded. If there is new firmware to be downloaded, it gets the data and writes it into program memory. There are many variations and additional features that can be added to improve reliability and simplify the use of the bootloader, some of which are discussed in this application note.

The general operation of a bootloader is discussed in the OPERATION section. Appendix A contains assembly code for a bootloader developed for the PIC16F877 and key aspects of this bootloader are described in the IMPLEMENTATION section.

For the purpose of this application note, the term "boot code" refers to the bootloader code that remains permanently in the microcontroller and the term "user code" refers to the user's firmware written into FLASH memory by the boot code.

FEATURES

The more common features a bootloader may have are listed below:

- Code at the Reset location.
- Code elsewhere in a small area of memory.
- Checks to see if the user wants new user code to be loaded.
- Starts execution of the user code if no new user code is to be loaded.
- Receives new user code via a communication channel if code is to be loaded.
- · Programs the new user code into memory.

OPERATION

The boot code begins by checking to see if there is new user code to be downloaded. If not, it starts running the existing user code. If there is new user code to be downloaded, the boot code receives and writes the data into program memory. There are many ways that this can be done, as well as many ways to ensure reliability and ease of use.

Integrating User Code and Boot Code

The boot code almost always uses the Reset location and some additional program memory. It is a simple piece of code that does not need to use interrupts; therefore, the user code can use the normal interrupt vector at 0×0004 . The boot code must avoid using the interrupt vector, so it should have a program branch in the address range 0×0000 to 0×0003 .

The boot code must be programmed into memory using conventional programming techniques, and the configuration bits must be programmed at this time. The boot code is unable to access the configuration bits, since they are not mapped into the program memory space. Setting the configuration bits is discussed in the next section.

In order for the boot code to begin executing the user code, it must know where the code starts. Since the boot code starts at the Reset vector, the user code cannot start at this location. There are two methods for placing the starting point of the user code.

One method is to use an ORG directive to force the user code to start at a known location, other than the Reset vector. To start executing the user code, the boot code must branch to this fixed location, and the user code must always use this same location as its start address.

An alternative method is to start the user code at the normal Reset vector and require that the user code has a goto instruction in the first four instructions to avoid the interrupt vector. These four instructions can then be relocated by the boot code and programmed into the area of program memory used by the boot code. This simplifies the development of code for use with the bootloader, since the user code will run when programmed directly into the chip without the boot code present. The boot code must take care of paging and banking so the normal Reset conditions apply before executing the relocated code.

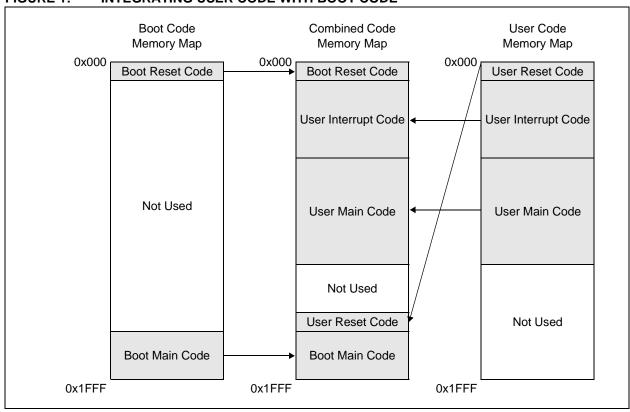


FIGURE 1: INTEGRATING USER CODE WITH BOOT CODE

Configuration Bits

The configuration bits cannot be changed by the boot code since they are not mapped into the program memory space. This means that the following configuration options must be set at the time that the boot code is programmed into the device and cannot be changed:

CPx Program Memory Code Protection Enable

DEBUG In-Circuit Debugger Mode Enable

WRT Program Memory Write Enable

CPD Data EEPROM Code Protection Enable

- LVP Low Voltage In-Circuit Programming Enable
- BODEN Brown-out Reset Enable
- PWRTE Power-up Timer Enable
- WDTE Watchdog Timer Enable
- FOSCx Oscillator Selection

Most of these configuration options are hardware or design-dependent, and being unable to change them when the user code changes is of no consequence.

The various PIC16F87X devices have different code protection implementations. Please consult the appropriate data sheet for details.

Some devices (such as the PIC16F877), can code protect part of the program memory and prevent internal writes to this protected section of memory. This can be used to protect the boot code from being overwritten, but also prevents the user code from being code protected, however.

On some devices, code protecting all the program memory still allows internal program memory write cycles. This provides security against the user code being read out of the chip, but does not allow the boot code to be protected from being overwritten.

Data EEPROM Code Protection Enable would normally not need to be set, unless data is programmed into the data EEPROM when the boot code is originally programmed and this data needs to be protected from being overwritten by the user code.

Program Memory Write Enable must be enabled for the boot code to work, since it writes to program memory. Low Voltage In-Circuit Serial Programming (ICSPTM) enable only needs to be set if the user wishes to program the PICmicro MCU in-circuit, using logic level signals on the RB3, RB6 and RB7 pins. Since the purpose of the boot code is to program user code into the PICmicro MCU, in most cases, it would be redundant to have facilities for low voltage ICSP.

If the Watchdog Timer is enabled, then the boot code must be written to support the Watchdog Timer and all user code will have to support the Watchdog Timer.

Determining Whether to Load New Code or to Execute User Code

After a Reset, the boot code must determine whether to download new user code. If no download is required, the bootcode must start execution of existing user code, if available.

There are many ways to indicate whether or not new user code should be downloaded. For example, by testing a jumper or switch on a port pin, polling the serial port for a particular character sequence, or reading an address on the l^2C^{TM} bus. The particular method chosen depends on the way that user code is transferred into the microcontroller. For example, if the new user code is stored on an l^2C EEPROM that is placed in a socket on the board, then an address in the EEPROM could be read to determine whether a new EEPROM is present.

If an error occurred while downloading new user code, or the bootloader is being used for the first time, there might not be valid user code programmed into the microcontroller. The boot code should not allow faulty user code to start executing, because unpredictable results could occur.

Receiving New User Code to Load into Program Memory

There are many ways that the microcontroller can receive the new firmware to be written into program memory. A few examples are from a PC over a serial port, from a serial EEPROM over an I^2C or SPITM bus, or from another microcontroller through the parallel slave port.

The boot code must be able to control the reception of data, since it cannot process any data sent to it while it is writing to its own program memory. In the case of data being received via RS-232, there must be some form of flow control to avoid data loss.

The data received by the boot code will usually contain more than just program memory data. It will normally contain the address to which the data is to be written and perhaps a checksum to detect errors. The boot code must decode, verify and store the data, before writing it into program memory. The available RAM (GPR registers) of the device limits the amount of data that can be received before writing it to program memory.

Programming the FLASH Program Memory

The PIC16F87X devices have special function registers that are used to write data to program memory. There is a specific sequence of writes to these registers that must be followed to reduce the chances of an unintended program memory write cycle occurring. Because code cannot be executed from the FLASH program memory while it is being written, program execution halts for the duration of the write cycle. Program memory is written one word at a time.

Error Handling

There are several things that can go wrong during execution of the boot code or user code. The bootloader should handle the following error conditions:

- No valid user code written into the chip.
- Error in incoming data.
- Received user code does not have any code at its Reset vector.
- Received user code overlaps boot code.
- User code causes execution into the boot code area.

If the bootloader is being used for the first time, or if the user code is partially programmed because of a previous error, there might not be valid user code programmed into the microcontroller. The boot code should not allow potentially faulty user code to start executing.

The transfer of data can be interrupted, which will cause the boot code to stop receiving data. There are several ways to handle this depending on how the data is being received. For example, the boot code may be able to time-out and request the data to be sent again. The simplest method is to wait, trying to receive more data with no time-out, until the user intervenes and resets the device. Since the boot code needs to leave the most possible program memory space for the user code and also be reliable, the smallest, simplest implementation is often the best.

Incoming data may be corrupted by noise or some other temporary interruption, and this should be detected, otherwise, incorrect data could be programmed. A checksum or other error detection method can be used.

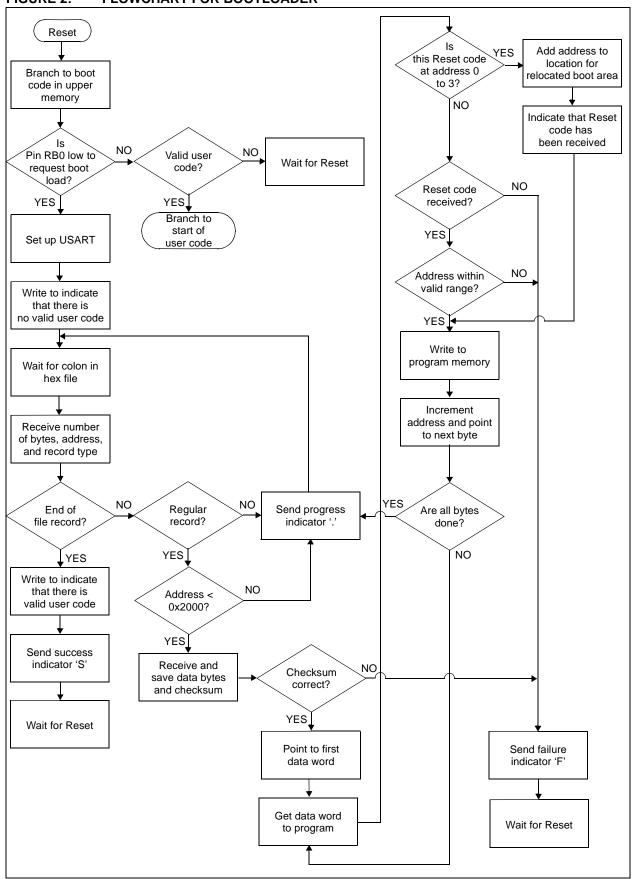
Incorrect use of flow control can result in data being sent to the PICmicro MCU while it is not ready to receive data. This can cause overrun errors that should be handled by the boot code. Once an overrun has occurred, the data is lost and this is essentially the same as a data transfer interruption, discussed above.

In some cases, data could be sent to the microcontroller before the boot code is running, causing part of the data to be lost. If this type of error is possible, then it should be detected. This error may manifest itself as user code that does not seem to have any code at the Reset location and can be detected by checking the addresses being programmed. An alternative is to generate a checksum on all the code that is written into program memory and transmit this to the user for verification, after programming has been completed. The code developer should take care that the user code does not use the same program memory space that the boot code uses. The exception is the user code at the Reset location that can be relocated, as explained earlier. If the user code does try to use program memory that contains boot code, the boot code should detect the conflicting address and not overwrite itself. In some devices, part of the program memory can be code protected to prevent internal writes to the part of the memory that contains the main boot code. Note that this does not apply to all PIC16F87X devices.

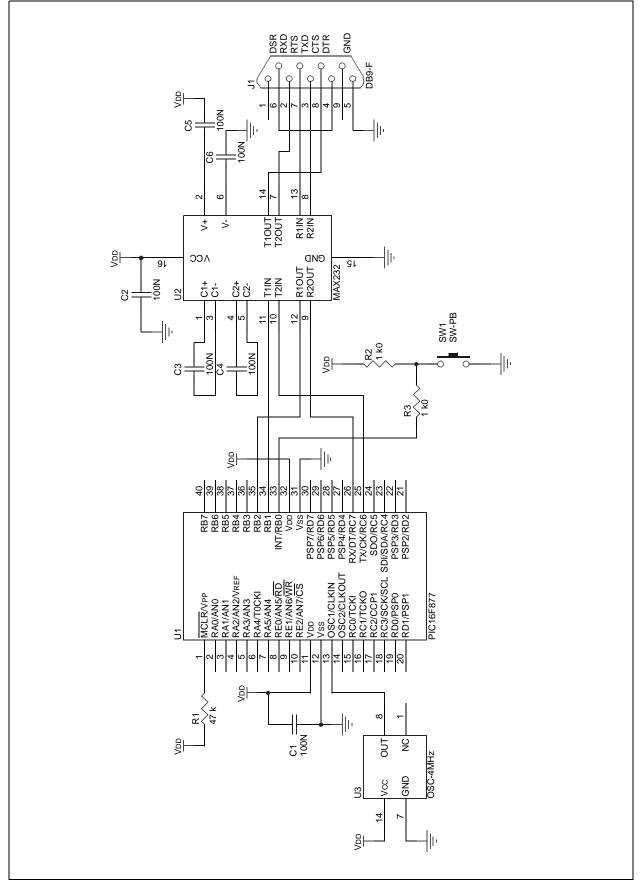
Faulty user code, or a brown-out condition that corrupts the program counter, can cause execution to jump to an unprogrammed memory location and possibly run into the start of the boot code. If the user code at the Reset location is being relocated, as explained earlier, then execution can enter the boot code area if a program branch does not occur in these four relocated instructions. The boot code should trap the program execution to avoid these errors from causing any unintended operation.

When an error is detected, it is useful to indicate this in some way. This can be as simple as turning on an LED, or sending a byte out the serial port. If the system includes a display and the display drivers are incorporated into the boot code, then more sophisticated error messages can be used.









IMPLEMENTATION

How this Bootloader Works

The boot code in Appendix A implements a bootloader in a PIC16F877 device. It uses the USART to receive data with hardware handshaking, tests a pin to decide if new user code should be received and includes many of the features discussed in this application note.

Integrating User Code and Boot Code

The code at the Reset location (ResetVector) writes to PCLATH. To set the page bits, it then jumps to the rest of the boot code in upper memory. The main code is in the upper 224 bytes of memory starting at address 0x1F20 (StartOfBoot). The first instructions at this location trap accidental entry into the boot code. The main bootloader routine starts at the address labeled Main.

The boot code requires that the user code includes a goto instruction in the first four locations after the Reset vector and relocates these four instructions into the boot code section (StartUserCode). This simplifies the development of code for use with the bootloader, since the same user code will also run when programmed directly into the chip, without the boot code present. The boot code changes to bank 0 and clears PCLATH before executing the relocated code, so that the normal Reset conditions apply. If a program branch does not occur in the four relocated instructions, then program execution is trapped in an endless loop to avoid any unintended operation.

The boot code must be programmed into the PIC16F877 using conventional programming techniques and the configuration bits are programmed at the same time. The configuration bits are defined with a _____CONFIG directive and cannot be accessed by the boot code, because they are not mapped into the program memory space. The boot code does not use a Watchdog Timer.

Determining Whether to Load new Code or to Execute User Code

The boot code tests port pin RB0 to determine whether new user code should be downloaded. If a download is required, then the boot code branches to the Loader routine that receives the data and writes it into program memory.

If pin RB0 does not indicate that new user code should be loaded, then a program memory location (labeled CodeStatus) is read with routine FlashRead to determine whether there is valid user code in the device. If there is valid user code, the boot code transfers execution to the user code by branching to location StartUserCode. Otherwise, execution is trapped in an endless loop to avoid this error from causing any unintended operation.

Receiving New User Code to Load into Program Memory

The boot code receives the new firmware as a standard Intel[®] hex file (INHX8M format), using the USART in Asynchronous Receiver mode (hex format defined in Appendix B). It is assumed that a PC will be used to send this file via an RS-232 cable, connected to a COM port. Hardware handshaking allows the boot code to stop the PC from transmitting data while FLASH program memory is being written. Since the PICmicro device halts program execution while the FLASH write occurs, it cannot read data from the USART during this time.

Hardware handshaking (described in Appendix C) is implemented using port pin RB1 as the RTS output and RB2 as the CTS input. The USART is set to 8-bit Asynchronous mode at 9600 baud in the SerialSetup routine. The SerialReceive routine enables transmission with the RTS output and waits until a data byte has been received by the USART, before returning with the data. The SerialTransmit routine checks the CTS input until a transmission is allowed and then sends a byte out the USART. This is used for transmitting progress indication data back to the PC.

The boot code receives the hex file, one line at a time and stops transmission after receiving each line, while received data is programmed into program memory.

Decoding the Hex File

The boot code remains in a loop, waiting until a colon is received. This is the first character of a line of the hex file. The following four pairs of characters are received and converted into bytes, by calling the GetHexByte routine. The number of bytes (divided by two to get the number of words) and the address (divided by two to get a word address) are saved, and the record type is checked for a data record, or end of file record.

If the record type shows that the line contains program memory data, then this data is received, two pairs of characters at a time (using the GetHexByte routine), and is stored in an array. The checksum at the end of the line is received and checked, to verify that there were not any errors in the line.

Once the hex file line has been received, hardware handshaking is used to stop further transmission, while the data is written into the program memory. The <CR> and <LF> characters that are sent at the end of the line are ignored. This gives the handshaking time to take effect by ignoring the byte being transmitted, when the handshaking signal is asserted. Once the data from the line has been programmed, the following lines are received and programmed in the same way, until the line indicating the end of the file has been received. A success indication 'S' is then transmitted out the USART (by the FileEnd routine) and the boot code waits for a Reset.

Programming the FLASH Program Memory

Data is written to the FLASH program memory using special function registers. The address is written to the EEADR and EEADRH registers and the first two bytes of data are written to EEDATA and EEDATH. The FlashWrite routine is then executed, which writes the data into program memory. The address is then incremented and the next two data bytes are written. This is repeated until all the data from the line of the hex file has been programmed into the FLASH program memory.

Error Handling

There are several things that can go wrong during execution of the boot code or user code, and a number of these error conditions are handled by the boot code. If an error occurs, the boot code traps it by executing an infinite loop, until the user intervenes and resets the device. If an error is detected in the incoming data, then a failure indication 'F' is transmitted. This does not occur in the case of an overflow error, or if the data transmission is halted.

If the bootloader is being used for the first time, or if the user code is partially programmed because of a previous error, there might not be valid user code programmed into the microcontroller. The boot code handles this by writing a status word (0x3fff) at a location labeled CodeStatus, before programming the FLASH device, and then writing a different status word (0x0000) to this same location, when programming of the user code has been completed. The boot code tests this location and only starts execution of the user code, if it sees that the user code was successfully programmed. When the boot code is originally programmed into the PICmicro MCU, the status word indicates that there is not valid user code in the device.

The transfer of data can be interrupted. In this case, the boot code waits, trying to receive more data with no time-out, until the user intervenes and resets the device. Noise, or a temporary interruption, may corrupt incoming data. The Intel hex file includes a checksum on each line and the boot code checks the validity of each line by verifying the checksum.

Incorrect use of flow control can result in data being sent to the PIC16F877, while it is not ready to receive data. This can cause an overrun error in the USART. Once an overrun has occurred, the USART will not move any new data into the receive FIFO and the boot code will be stuck in a loop waiting for more data. This effectively traps the error until the user intervenes by resetting the device.

If the user starts transmitting a hex file before the boot code is running, the boot code may miss the first lines of the file. Since all the lines of a hex file have the same format, it is not normally possible to determine whether the line being received is the first line of the hex file. However, since MPASM generates hex files with addresses in ascending order, the first valid line of the hex file should contain the code for the Reset vector which is checked by the boot code.

The user code may try to use program memory locations that contain boot code. This is detected by checking the address being programmed and detecting conflicting addresses. The boot code will not overwrite itself and is not code protected.

Faulty user code, or noise that corrupts the program counter, can cause execution to jump to an unprogrammed memory location and possibly run into the start of the boot code. The first instructions in the boot code are an infinite loop that traps execution into the boot code area.

Because the first four instructions in program memory are relocated in the boot code implementation, there must be a program branch within these four instructions. If there is no program branch, then execution is trapped by the boot code.

Using the Bootloader

The procedure for using the bootloader is as follows:

- On the PC, set up the serial port baud rate and flow control (hardware handshaking).
- Connect the serial port of the PIC16F87X device to the serial port of the PC.
- Press the switch to pull pin RB0 low.
- Power up the board to start the boot code running.
- The switch on RB0 can be released if desired.
- From the PC, send the hex file to the serial port.
- A period '.' will be received from the serial port for each line of the hex file that is sent.
- An 'S' or 'F' will be received to indicate success or failure.
- The user must handle a failure by resetting the board and starting over.
- Release the switch to set pin RB0 high.
- Power-down the board and power it up to start the user code running.

On the PC, there are several ways to set up the serial port and to transfer data. This also differs between operating systems.

A terminal program allows the user to set up and send data to a serial port. In most terminal programs, an ASCII or text file can be sent and this option should be used to send the hex file. A terminal program will also show data received on the serial port and this allows the user to see the progress '.' indicators and the success 'S' or failure 'F' indicators. There are many terminal programs available, some of which are available free on the Internet. This boot code was tested using Tera Term Pro, Version 2.3. The user should be aware that some popular terminal programs contain bugs. A serial port can be set up in a DOS window, using the MODE command and a file can be copied to a serial port, using the COPY command. When using Windows[®] 95/98, the MODE command does not allow the handshaking signals to be configured. This makes it difficult to use the COM port in DOS. When using Windows NT[®] or Windows 2000[®], the following commands can be used to send a hex file named filename.hex to serial port COM1:

MODE COM1: BAUD=9600 PARITY=N DATA=8 STOP=1 to=off xon=off odsr=off octs=on dtr=off rts=on idsr=off

COPY filename.hex COM1:

Resources Used

The boot code coexists with the user code on the PIC16F877 and many of the resources used by the boot code can also be used by the user code. The boot code uses the resources listed in Table 1.

TABLE 1: RESOURCES USED BY THE BOOT CODE

Resource	Amount
Program memory	224 words
Data memory	72 bytes
I/O pins	5 pins
Peripherals	USART

The program memory used by the boot code cannot be used for user code, although it is possible to call some of the subroutines implemented in the boot code to save space. The user code can use all the data memory.

The USART can be used by the user code with the two I/O pins for the USART and the I/O pins used for handshaking. The I/O pin used to indicate that the boot code should load new user code, is connected to a switch or jumper. This can be isolated with a resistor and used as an output, so that it is possible to use all the I/O pins used by the bootloader.

In summary, all resources used by the boot code, except program memory, can also be used by the user code.

CONCLUSION

Using a bootloader is an efficient way to allow firmware upgrades in the field. Less than 3% of the total program memory is used by the boot code and the entire program memory available on a PIC16F877 can be programmed in less than one minute at 19,200 baud.

The cost of fixing code bugs can be reduced with a bootloader. Products can be upgraded with new features in the field, adding value and flexibility to the products. The ability to upgrade in the field is an added feature and can enhance the value of a product.

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00000027	00068	DataPointer		lata in buffer
0000028	00069	DataArray		;buffer for storing incoming data
	000700071	ENDC		
	00073	; to select the	register bank	
	00074	;Many bank changes can		when only one STATUS bit changes
	00076	Bank0 MACRO		;macro to select data RAM bank 0
	00077		STATUS, RP0	
	00078	bcf	STATUS, RP1	
	00079	ENDM		
	00081	Bank1 MACRO		;macro to select data RAM bank 1
	00082		STATUS, RPO	
	00083	bcf	STATUS , RP1	
	00084	ENDM		
	00085			
	00086	Bank2 MACRO		;macro to select data RAM bank 2
	00087	bcf	STATUS, RPO	
	00088	bsf	STATUS, RP1	
	00089	ENDM		
	06000			
	0000	Bank3 MACRO		;macro to select data RAM bank 3
	00092	bsf	STATUS, RPO	
	00033	bsf	STATUS, RP1	
	00094	ENDM		
	00095			
	0000	======================:=:		
	00097	;Reset vector code		
0	00098			
0000	66000	OKG	nxnnn	
4TOS 0000	10100	Kesetvector: movim	nign Main	10000 1110 1000
LOOD TOUL				SEC PAGE NICS FOF PAGES
Message[306]: 0000 2500	Crossing page	boundary ensu	e bits are set. Min	
	00104	90-00 90-00	MALII	
	00105			
	00100	;Start of boot code in	. upper memory tra	traps accidental entry into boot code area
	00102			
1F20	00108	ORG	0x1f20	;Use last part of page3 for PIC16F876/7
	00100	; ORG	0×0f20	оf
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<pre>;trap if execution runs into boot code ;set correct page ;trap error and wait for reset</pre>	p to start of user code to this routine	;set correct page for reset condition	; relocated user code replaces this nop	; relocated user code replaces this nop	user code	; relocated user code replaces this nop	;trap if no goto in user reset code	;set correct page	trap error and wait for reset;		r valid code has been programmed		;0 for valid code, 0x3fff for no code				d if valid user code exists		change	T ;check pin for boot load	; if low then do bootload	; load address of CodeStatus word	; read data at CodeStatus location		;set Z flag if data is zero		t Z flag	; if not zero then is no valid code	; if zero then run user code			ew code		; indicate no reset vector code yet		; load address of CodeStatus word	;load data to indicate no program		;load data to indicate no program
high TrapError PCLATH TrapError	code to jump to survey to this to the second s	PCLATH					high TrapError1	PCLATH	TrapError1		n to show whether valid		0x3fff				I should occur and			PORTB, TEST_INPUT	Loader	LoadStatusAddr	FlashRead		EEDATA, F		STATUS, Z	TrapError2	StartUserCode			load and program new		TestByte		LoadStatusAddr	0x3f	EEDATH	0xff
movlw movwf goto	reset c ko befor	clrf	dou	dou	dou	dou	movlw	movwf	goto		locatic		DA			code routine	if a load		Bank0	btfss	goto	call	call	Bank2	movf	BankO	btfss	goto	goto			t t		clrf		call	movlw	movwf	movlw
StartOfBoot: TrapError:	;Relocated user reset code to jum ;Must be in bank0 before jumping	StartUserCode:							TrapError1:		Program memory location to		CodeStatus:		:	; Main boot code	Tests to see i		Main:									TrapError2:			:	;Start of routine		Loader:					
00112 00113 00114 00115	00117 00118 00118	00120	00121	00122	00123	00124	00125	00126		00128	00130	00131	00132	00133	00134	00135	00136	00137	00138	00139	00140	00141	00142	00143	00144	00145	00146	00147	00148	00149	00150	00151	00152	00153	00154	00155	00156	00157	00158
301F 008A 2F22		018A	0000	0000	0000	0000	301F	008A	2F2A				3 F F F							1C06	2F3A	27B5	27F6		088C		1D03	2F38	2F23					01A5		27B5	303F	008E	30FF
1F20 1F21 1F22		1F23	1F24	1F25	1F26	1F27	1F28	1F29	1F2A				1F2B							1F2E	1F2F	1F30	1F31		1F34		1F37	1F38	1F39					1F3A		1F3B	1F3C	1F3D	1F3E

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;write new CodeStatus word	;set up serial port		1	t ':' t address and number of bytes	1	;get new byte from serial port	; check if ':' received		; if not then wait for next byte		;start with checksum zero			; limit number in case of error in file			;divide by 2 to get number of words		;get upper half of program start address			;get lower half of program start address				;divide address by 2 to get word address			;get record type		; check if end of file record (0x01)	; if end of file then all done				;check if regular line record (0x00)	; if not then ignore line and send '.'			;check if address < 0x2000	; which is ID locations and config bits	then iqnore line and)
EEDATA FlashWrite	SerialSetup			<pre>tile starting with ter ':' and extract</pre>		SerialReceive	,:,	STATUS, Z	GetNewLine		Checksum		GetHexByte	0×1F	NumWords	STATUS, C	NumWords, F		GetHexByte	AddressH		GetHexByte	AddressL		STATUS, C	AddressH, F	AddressL, F		GetHexByte	0×01	STATUS, Z	FileDone		HexByte,W	0×00	STATUS, Z	LineDone		0xe0	AddressH,W	STATUS, C	LineDone	
movwf call	call			ot hex oytes af	I	call	xorlw	btfss	goto		clrf		call	andlw	movwf	bcf	rrf		call	movwf		call	movwf		bcf	rrf	rrf		call	xorlw	btfsc	goto		movf	xorlw	btfss	goto		movlw	addwf	btfsc	goto	1
00159 00160	00161 00162	00163		00165 ;Get new line 00166 ;Get first 8]	00167	00168 GetNewLine:	00169	00170	00171	00172	00173	00174	00175	00176	00177	00178	00179	00180	00181	00182	00183	00184	00185	00186	00187	00188	00189	00190	00191	00192	00193	00194	00195	00196	00197	00198	00199	00200	00201	00202	00203	00204	00205
					0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
008C 27EA	27CC					2 7 D B	3A3A	1D03	2F42		01A3		2 7BC	391F	0 0A2	1003	0 CA2		27BC	0 0 A 0		27BC	0 0A 1		1003	0 CA 0	0 CA 1		27BC	3A01	1903	2 FAB		0826	3A00	1D03	2 FA 8		30E0	0720	1803	2 FA 8	
1F3F 1F40	1F41					1F42	1F43	1F44	1F45		1F46		1F47	1F48	1F49	1F4A	1F4B		1F4C	1F4D		1F4E	1F4F		1F50	1F51	1F52		1F53	1F54	1F55	1F56		1F57	1F58	1F59	1F5A		1F5B	1F5C	1F5D	1F5E	

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1F5F 3(1F60 0(1F61 08 1F62 0(1F63 2) 1F63 2) 1F65 0(1F65 02	3 3 2 2 8 0 0 8 4 0 0 2 4 0 0 2 4 2 7 B C 0 0 2 8 4 0 2 8 4	00207; 00208 00209 00210 00211 00213 00213 00215 00215	;Get data bytes GetData:		and checksum from line movlw DataArray movwf FSR movf NumWords,W movwf Counter call GetHexByte movwf INDF incf FSR,F	of hex file ;set pointer to start of array ;set counter to number of words ;get low data byte ;save in array ;point to high byte
1F66 2' 1F67 0(1F68 0/ 1F69 0 1F6A 2] 1F6A 2] 1F6B 2'	27BC 080 0A84 0BA4 2F63 27BC	00219 00219 00220 00221 00221 00222 00223 00223		call movwf incf decfsz goto call	GetHexByte INDF FSR,F Counter,F GetData GetHexByte	;get high data byte ;save in array ;point to next low byte ;get checksum
1F6C 08 1F6D 1D 1F6E 2F 1F6F 14 1F6F 14	0823 1D03 2FB2 1486	00226 00227 00228 00229 00230 00231 00233 700233 700233		movf (btfss (goto H bsf H	Checksum, W STATUS, Z BrrorMessage PORTB, RTS_OUTF 	<pre>;check if checksum correct UT ;set RTS off to stop data being received</pre>
1F70 3(1F71 0(1F72 08 1F73 0(1F73 0(3 0 2 8 0 0 8 4 0 0 8 2 2 0 0 A 4 0 0 A 4		m mc mc mc check if address	movlw movwf movf movwf sis in	DataArray FSR NumWords,W Counter 	point to start of array set counter to half number of bytes
1F74 08 1F75 11 1F75 21 1F77 30 1F77 30 1F77 21 1F73 21 1F73 21 1F73 21	0820 11003 2F84 3 30FC 0721 1803 2F84 2F84		CheckAddress:	movf btfss goto movlw btfsc goto	AddressH,W STATUS,Z CheckAddress1 Oxfc AddressL,W STATUS,C CheckAddress1	<pre>;checking for boot location code ;test if AddressH is zero ;if not go check if reset code received ;add 0xfc (-4) to address ;no carry means address < 4 ;if not go check if reset code received</pre>

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AddressL,W ;relocate addresses 0-3 to new location low (StartUserCode + 1) :add low address to new location	from bank0 to bank2	ew low address	(StartUserCode + 1) ;get new location high address	lress	;go get data byte and program into flash			red	ייטים בייויי דד הייביי בייני	VIITTICLE MICH DOOL TOARET		; check 1	; if not then error		Boot ;get high byte of address		test if less than boot code address;	so continue with write		, and an arrive in high hite of addree	THO SO STICT IN WIGH DACE OF	. 200 + 1011 - 1011 - 101	; der tow byre		;test if less than boot code address	e ;no so error in address			ia into flash		;get high address	change from bank0 to bank2	;load high address	; change from bank2 to bank0	;get low address	; change from bank0 to bank2	; load low address		get low byte from array;	low byte	point to high data byte	bvte from		. noint to next low data byte		.write data to program memory	WITC CALA IO NTOTAM MCMOTA	
AddressL,W low (StartUs		EEADR	high (StartU	EEADRH	LoadData			has been received	ל י- ל	116111		Testbyre, u	ErrorMessage		high StartOfBoot	AddressH,W	STATUS, C	LoadAddress	STATIS Z	ErrorMeesence	DERCENTOTIO		TOW PLAFLUIDOUL	AddressL,W	STATUS, C	ErrorMessage			and write data		AddressH,W		EEADRH		AddressL,W		EEADR		INDF,W	EEDATA	FSR, F	INDF.W	EEDATH	F.S.R. F	+ / / +	FlachWrite	L TODIME TOO	
movf addlw	Bank2	movwf	movlw	movwf	goto			Check if reset code he		AT AUDIDOD IT		UneckAddress1: DTISS	goto		movlw	subwf	btfss	aoto	bt fas		<u> </u>	•• [••····		SUDWI	btfsc	goto			;Load address and data		LoadAddress: movf	Bank2	movwf	BankO	movf	Bank2	movwf		LoadData: movf	movwf	incf	movf	movwf	incf		רפט	1 1 2 2	
00253	00255	00256	00257	00258	00259	00260	00261						00266	00267	00268	00269	00270	00271	00272	2/200 6/000			G/700	00276	00277	00278	00279	00280		00282		00284	00285	00286	00287	00288	00289	00290	00291	00292	00293	00294	00295	00296	00297	00098	00299	
0821 3E24) 1 1	008D	301F	008F	2F9A						L C T	TCZ 2	2FB2		301F	0220	1C03	2F90	1 D.0.3		21D2		0202	0221	1803	2FB2					0820		008F		0821		008D		0800	0080	0A84	0800	008E	0284		0 7 F D	4	
1F7C	i	1F80	1F81	1F82	1F83							LF'84	1F85		1F86	1F87	1F88	1F89	1 F 8 A		T O D			1F8D	1F8E	1F8F					1F90		1F93		1F96		1F99		1F9A	1F9B	1F9C	1F9D	1F9E	1 F 9 F	1 1 4 1	1 600		

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.cat new hyte from serial nort	, שבר זופא הערב דדטוו אבודמד הטור	;add -'A' to Ascii high byte	;check if positive	;if not, add 17 ('0' to '9')	;else add 10 ('A' to 'F')	save nibble:		INDAE ITTATA LO ITATI POSTLITATI		new byte from	;add -'A' to Ascii low byte	;check if positive	;if not, add 17 ('0' to '9')	;else add 10 ('A' to 'F')	; add low nibble to high nibble	put result in W reg	add to cumulative checksum					placed in-line saving a call and return			;change from bank3 to bank0	PORTB, RTS_OUTPUT ;set RTS off before setting as output	; change from bank0 to bank1	TRISB,RTS_OUTPUT ;enable RTS pin as output	set baud rate 9600 for 4Mhz clock		.hand rate high sneed ontion				; enable reception	; enable serial port				for bute to be received in IISDPT and return with bute in W				; change from unknown bank to bank0	';set RTS on for data to be received	check if data received		; set received data into W
סאייס <i>ר</i> סק רביידסט	A TATVACTA	0xbf	STATUS, C	0×07	0x0a	HexBvte		лехьуге, г	-	SerialReceive	0xbf	STATUS, C	0×07	0x0a	HexByte, F	HexByte, W	Checksum, F					nd can be	n bank0			PORTB, RTS_OUTPUT		TRISB, RTS_OUTPUT	BAUD CONSTANT	SPBRG	нряя атехт	TYSTA TYFN	MENT VETONI		RCSTA, CREN	RCSTA, SPEN				a TISART a better	CETVED TH DOWN W	n bank0			PORTB, RTS_OUTPUT	PIR1, RCIF	ъ-1 -	RCREG, W
ר נ	COLL	addlw	btfss	addlw	addlw	movwf	f down	SWAPL	,	call	addlw	btfss	addlw	addlw	iorwf	movf	addwf	return			or async	Y called	ceturns i		BankO	bsf	Bank1	bcf	movlw	movwf	haf	а с ц ц ц ц		Banku	bsf	bsf	return			to he re	י ש ר ר ר ר ר	ceturns 1		Bank0	bcf	btfss	aoto	movf
Санцахвита.	se unexpy re:																				;set up USAKI IOF	koutine is only called once and can	;This routine returns in bank0		SerialSetup:															·Wait for hyte		This routine returns in bank0;		SerialReceive:				
00348		00350	00351	00352	00353	00354				00357	00358	00359	00360	00361	00362	00363	00364	00365	00366	00367				00371	00372 5	00373	00374	00375	00376	00377	87200	0/200		00380	00381	00382	00383	00384	00385				00388	00389	00390	00391	00392	00393
aur c	au 1 2	3 EBF	1C03	3 E 0 7	3E0A	0046		OFAO		2 7DB	3 EBF	1C03	3 E 0 7	3E0A	04A6	0826	0 7 A 3	0008								1486		1086	3019	6600	1518	1698	2		1618	1798	0008								1086	1E8C	2 FDE	081A
С Ц Ц Ц		1 FBD	1FBE	1FBF	1FC0	1FC1		TLCZ		1FC3	1FC4	1FC5	1FC6	1FC7	1FC8	1FC9	1 F C A	1 FCB								1FCE		1FD1	1FD2	1FD3	1 FD4	- CH	н. Г Г		1FD8	1FD9	1 FDA								1 F D D	1 FDE	1 FDF	1750

1FE4 1906 1FE5 2FE4 1FE5 2FE4 1FE6 1E0C 1FE7 2FE6 1FE8 0099 1FE9 0008	00399 00399 00399 00399 00400 00400 00400 00400 000400 000400 000400 000400 000400 000400 000400 000400 000400 000400 000400 000400 000400 00000 00000 00000 00000 00000 00000 0000	<pre>;syte in W register from USART ;This routine returns in bank0 SerialTransmit: Bank0 btfsc PORTB,CTS_INPUT goto \$-1 btfss PIR1,TXIF goto \$-1 movwf TXREG return</pre>	jister from USART in banko PORTB,CTS_INPUT \$-1 PIR1,TXIF \$-1 TXREG	;change from unknown bank to bank0 ;check CTS to see if data can be sent ;check that buffer is empty ;transmit byte
1FEC 3084 1FED 008C 1FEE 3055 1FEF 008D 1FF1 008D 1FF1 008D 1FF2 148C	000410 00410 004110 004114 004114 004114 004114 004114 004114 004117 004119 004119 004210 004210	<pre>;Write to a location in the fl ;Address in EEADRH and EEADR, ;This routine returns in bank3 FlashWrite: Bank3 0x84 movuf EECON1 movuf EECON2 movuf EECON2 movuf EECON2 bsf EECON1</pre>	ash prog data in , WR	EEDATH and EEDATA EEDATH and EEDATA ;change from bank2 to bank3 ;enable writes to program flash ;do timed access writes ;begin writing to flash
1FF3 0000 1FF4 0000 1FF5 0008 1FF6 301F 1FF7 058F 1FFA 3080 1FFB 008C 1FFB 008C	00423 00425 00425 00425 00424 00423 00423 00433 00433 00433 00433 00433 00433 00433 00433 00433 00433 00433 00433 0000433 0000433 0000433 000433 000433 000433 000433 000433 0000433 000433 000433 000433 000433 000433 000433 000433 000433 000433 000433 000433 000433 000433 00000000	nop nop return ;	the flash produce and can be once and can be once and can be once and is oxif EEADRH, F EEADRH, F EECON1 PD	<pre>;processor halts here while writing </pre>

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processor waits while reading:					
nop ; processor	return	END			
41 42 42	44 45 46 2	*/ 48 sed, '-' = Unused)		nused. : 227 : 7965	0 suppressed
004410000044200000000000000000000000000		00448 00448 MEMORY USAGE MAP ('X' = Used,	XXXX XXXXXXXXXXXXXXXXXXXXXXXXXXXX	All other memory blocks unused. Program Memory Words Used: 22 Program Memory Words Free: 796	: 0 : 0 reported, : 1 reported,
1FFD 0000		MEMORY US.	0000 : XXX 1F00 : 1F40 : XXX 1F80 : XXX 1FC0 : XXX 2000 :	All other Program Me Program Me	Errors Warnings Messages

APPENDIX B: HEX FILE FORMAT

MPASM generates an 8-bit Intel hex file (INHX8M) by default. The lines of this hex file all have the following format:

: BBAAAATTHHHH.... HHCC

A colon precedes each line and is followed by hexadecimal digits in ASCII format.

BB is a 2-digit hexadecimal byte count representing the number of data bytes that will appear on the line. This is a number from 0x00 to 0x10 and is always even because the PIC16F87X parts have a 14-bit wide memory and use two bytes for every program memory word.

AAAA is a 4-digit hexadecimal address representing the starting byte address of the data bytes that follow. To get the actual program memory word address, the byte address must be divided by two.

TT is a 2-digit hexadecimal record type that indicates the meaning of the data on the line. It is 0×00 for a regular data record and 0×01 for an end of file record. The boot code ignores all other record types.

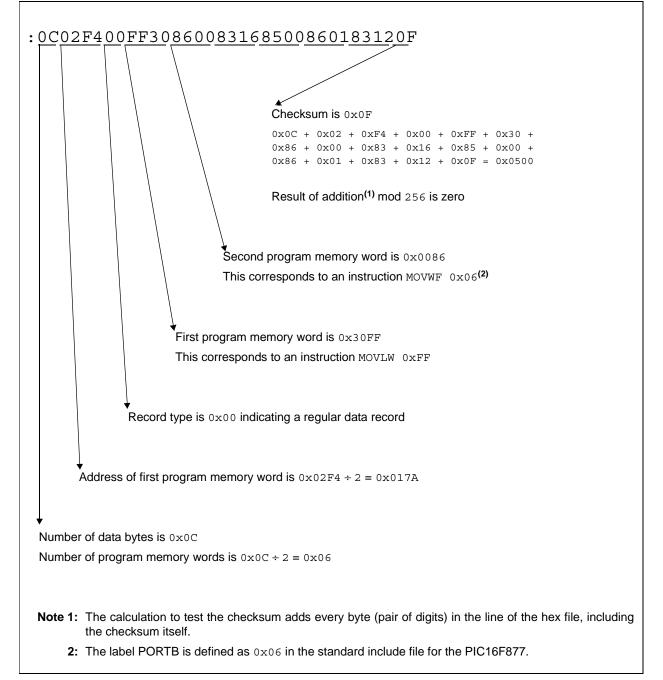
HH are 2-digit hexadecimal data bytes that correspond to addresses, incrementing sequentially from the starting address earlier in the line. These bytes come in low byte, high byte pairs, corresponding to each 14-bit program memory word.

cc is a 2-digit hexadecimal checksum byte, such that the sum of all bytes in the line including the checksum, is a multiple of 256. The initial colon is ignored. The code in Example B-1 will generate a line in a hex file as shown in Figure B-1.

EXAMPLE B-1: CODE TO GENERATE A HEX FILE

ORG	0x17A
movlw	0xFF
movwf	PORTB
bsf	STATUS, RPO
movwf	TRISA
clrf	TRISB
bcf	STATUS, RPO

FIGURE B-1: LINE OF HEX FILE



APPENDIX C: RS-232 HARDWARE HANDSHAKING SIGNALS

Understanding hardware flow control can be confusing, because of the terminology used and the slightly different way that handshaking is now implemented, compared to the original specification.

RS-232 hardware handshaking was specified in terms of communication between Data Terminal Equipment (DTE) and Data Communications Equipment (DCE). The DTE (e.g., computer terminal) was always faster than the DCE (e.g., modem) and could receive data without interruption. The hardware handshaking protocol required that the DTE would request to send data to the DCE (with the request to send RTS signal) and that the DCE would then indicate to the DTE that it was cleared to send data (with the clear to send CTS signal). Both RTS and CTS were, therefore, used to control data flow from the DTE to the DCE.

The Data Terminal Ready (DTR) signal was defined so that the DTE could indicate to the DCE that it was attached and ready to communicate. The Data Set Ready (DSR) signal was defined to enable the DCE to indicate to the DTE that it was attached and ready to communicate. These are higher level signals not generally used for byte by byte control of data flow, although they can be used for this purpose.

Most RS-232 connections use 9-pin DSUB connectors. A DTE uses a male connector and a DCE uses a female connector. The signal names are always in terms of the DTE, so the RTS pin on the female connector of the DCE is an input and is the RTS signal from the DTE. Over time, the clear distinction between the DTE and DCE has been lost. In many instances, two DTE devices are connected together. In other cases, the DCE device is able to send data at a rate that is too high for the DTE to receive continuously. In practice, the DTR output of the DTE has come to be used to control the flow of data to the DTE and now indicates that the DCE (or other DTE) may send data. It no longer indicates a request to send data to the DCE.

It is common for a DTE to be connected to another DTE (e.g., two computers), and in this case, they will both have male connectors and the cable between them will have two female connectors. This is known as a null modem cable. The cable is usually wired in such a way that each DTE looks like a DCE to the other DTE. To achieve this, the RTS output of one DTE is connected to the CTS input of the other DTE and vice versa. Each DTE device will use its RTS output to allow the other DTE device to transmit data and will check its CTS input to determine whether it is allowed to transmit data.

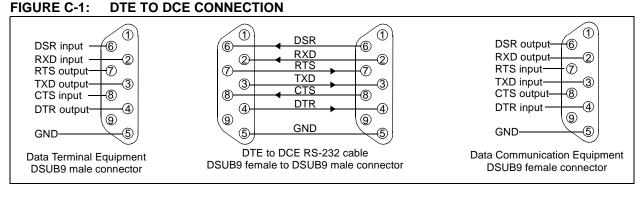
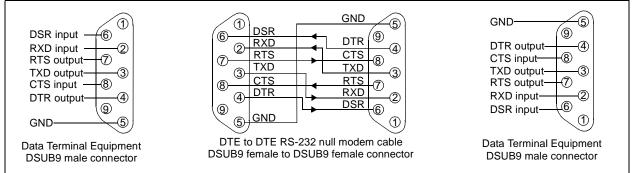


FIGURE C-2: DTE TO DTE CONNECTION



Note the following details of the code protection feature on PICmicro[®] MCUs.

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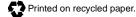
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